



## **Chap 6: Spatial Networks**

6.1 Example Network Databases
6.2 Conceptual, Logical and Physical Data Models
6.3 Query Language for Graphs
6.4 Graph Algorithms
6.5 Trands: Access Methods for Spatial Networks

**6.5 Trends: Access Methods for Spatial Networks** 



## **Learning Objectives**

### Learning Objectives (LO)

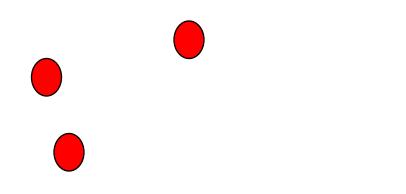
- LO1: Understand the concept of spatial network (SN)
  - What is a spatial network?
  - Why learn about spatial network?
- LO2 : Learn about data models for SN
- LO3: Learn about query languages and query processing
- LO4: Learn about trends
- Focus on concepts not procedures!
- Mapping Sections to learning objectives

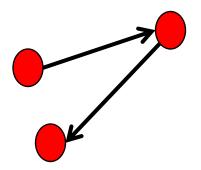
ф	LO1	-	6.1
ф	LO2	-	6.2
ቍ	LO3	-	6.3, 6.4
ф-	LO4	-	6.5





## From proximity to connectivity

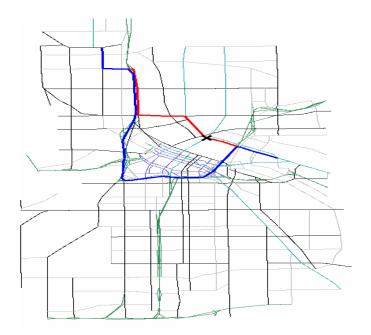






**6.1 Example Spatial Networks** 

#### **Road networks**

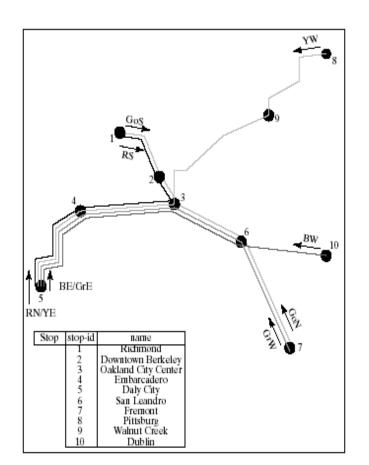


- **1.** Find shortest path from my current location to a destination.
- **2.** Find nearest hospital by distance along road networks.
- **3.** Find shortest route to deliver goods to a list of retail stores.
- **4.** Allocate customers to nearest service center using distance along roads



## **6.1 Example Spatial Networks**

#### **Railway networks**



- **1.** Find the number of stops on the Yellow West (YW) route.
- **2.** List all stops which can be reached from Downtown Berkeley.
- **3.** List the route numbers that connect Downtown Berkeley and Daly City.
- **4.** Find the last stop on the Blue West (BW) route.



## **6.1 Example Spatial Networks**

#### **River networks**



**1.** List the names of all direct and indirect tributaries of the Mississippi river

**2.** List the direct tributaries of the Colorado.

**3.** Which rivers could be affected if there is a spill in river P1?



## **Learning Objectives**

- Learning Objectives (LO)
  - LO1: Understand the concept of spatial network (SN)
  - LO2 : Learn about data models of SN
    - Representative data types and operations for SN
    - Representative data-structures
  - LO3: Learn about query languages and query processing
  - LO4: Learn about trends
- Focus on concepts not procedures!
- Mapping Sections to learning objectives

ф	LO1	-	6.1
ф	LO2	-	6.2
¢	LO3	-	6.3, 6.4

🛚 LO4 - 6.5



## **6.2 Spatial Network Data Models**

- •Recall 3 level Database Design
  - •Conceptual Data Model
    - Graphs
  - •Logical Data Model -
    - •Data types Graph, Vertex, Edge, Path, ...
    - •Operations Connected(), Shortest\_Path(), ...
  - •Physical Data Model
    - •Record and file representation adjacency list
    - •File-structures and access methods CCAM



## **6.2 Conceptual Data Models**

- Conceptual Data Model for Spatial Networks
  - A graph, G = (V,E)
  - V = a finite set of *vertices*
  - E = a set of edges E , between vertices in V

Classifying graph models

•Do nodes represent spatial points? - spatial vs. abstract graphs

•Are vertex-pair in an edge order? - directed vs. undirected

Examples

- Road network is a spatial graph, River network is an abstract graph
- River network is directed, Road network can be directed or undirected



## 6.2 Physical Data Models

- Categories of record/file representations
  - Main memory based
  - Disk based
- •Main memory representations of graphs
  - Adjacency matrix M[A, B] = 1 if and only if edge(vertex A, vertex B) exists
  - •Adjacency list : maps vertex A to a list of successors of A
  - •Example: See Figure 6.2(a), (b) and (c) on next slide

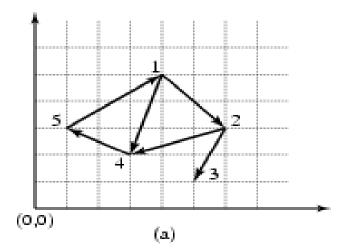
•Disk based

- •normalized tables, one for vertices, other for edges
- •denormalized one table for nodes with adjacency lists
- •Example: See Figure 6.2(a), (d) and (e) on next slide

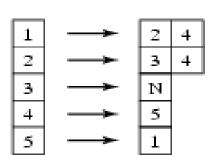


## **<u>6.2.2 Physical Data Models - Figure 6.2</u>**

Fig 6.2



		Destination				
		1	2	3	4	5
	1	0	1	0	1	0
Q	2	0	0	1	1	0
source	3	0	0	0	0	0
й	з 4	0	0	0	0	1
	5	1	0	0	0	0





(c) Adjacency-List

Node (R)			
id	x y		
1	4.0	5.0	
2	6.0	3.0	
3	5.0	1.0	
4	3.0	2.0	
5	1.0	3.0	

Edge (S)				
scurce	dest	distance		
1	2	,s		
1	4	√10		
2	3	√5		
2	4	√10		
4	5	√5		
5	1	$\sqrt{18}$		

(d) Node and Edge Relations

id	x	у	Successors	Predecessors
1	4.0	5.0	(2,4)	(5)
2	6.0	3.0	(3,4)	(1)
3	5.0	1.0	0	(2)
4	3.0	2.0	(5)	(1,2)
5	1.0	3.0	(1)	(4)

(e) Denormalized Node Table



## **Learning Objectives**

- Learning Objectives (LO)
  - LO1: Understand the concept of spatial network (SN)
  - LO2 : Learn data models for SN
  - LO3: Learn about query languages and query processing
    - Query building blocks
    - Processing strategies
  - LO4: Learn about trends
- Focus on concepts not procedures!
- Mapping Sections to learning objectives

ф	LO1	-	6.1
ф	LO2	-	6.2
ф	LO3	-	<mark>6.3,</mark> 6.4
			<b>a -</b>

🛚 LO4 - 6.5

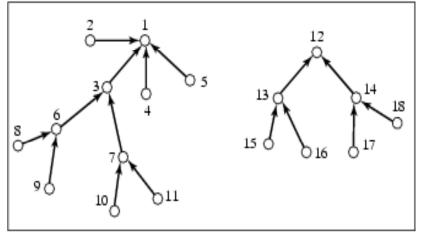


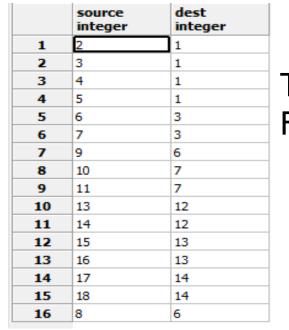
- Recall Relation algebra (RA) based languages
  - Can not compute paths with arbitrary length
  - •SQL support for graph queries
  - SQL2 CONNECT clause in SELECT statement
    - For directed acyclic graphs, e.g. hierarchies
  - SQL 3 WITH RECURSIVE statement
    - Transitive closure on general graphs
  - SQL 3 -user defined data types
    - Can include shortest path operation!



## 6.3.2 SQL2 Connect Clause

- Syntax details
  - FROM clause a table for directed edges of an acyclic graph
  - PRIOR identifies direction of traversal for the edge
  - START WITH specifies first vertex for path computations
- Semantics
  - List all nodes reachable from first vertex using directed edge in specified table
  - Assumption no cycle in the graph!
  - Not suitable for train networks, road networks





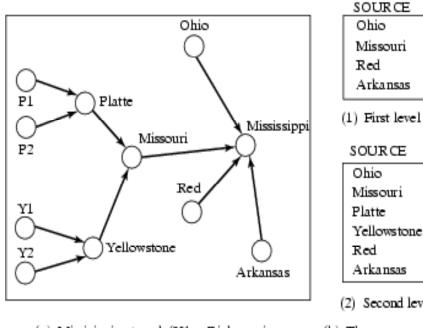
#### Table FallsInto

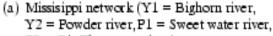


## SQL Connect Clause - Example

- SQL experssion on right
- Execution trace of paths
  - •starts at vertex 1 (Mississippi)
  - •adds children of 1
  - •adds children of Missouri
  - adds children of Platte
  - adds children of Yellostone
- Result has edges
  - •from descendents
  - to Mississippi

**SELECT** source **FROM** FallsInto **CONNECT BY PRIOR** source = dest START WITH dest =1





P2 = Big Thompson river).

Ohio Missouri Platte P1 **P**2 Yellowstone  $\mathbf{Y}1$ Y2Red Arkansas (3) Final result

SOURCE

(2) Second level

Arkansas

(b) The sequence of the CONNECT clause

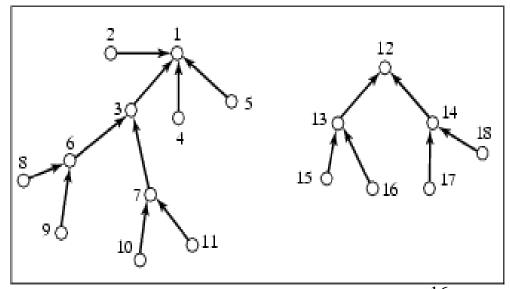


## **SQL Connect Clause - Exercise**

- Study 2 SQL queries on right
  - Note different use of PRIOR keyword
- SELECT source FROM FallsInto CONNECT BY PRIOR source = dest START WITH dest =3

- Compute results of each query
- Which one returns ancestors of 3?
- •Which returns descendents of 3?
- Which query lists river affected byoil spill in Missouri (id = 3)?

SELECT source FROM FallsInto CONNECT BY source = PRIOR dest START WITH dest =3

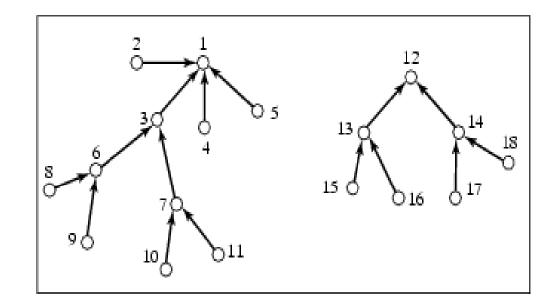




## **SQL Connect Clause - Exercise**

- Stop to a specified level
- •Visit 2, 3, 4, 5, 6, 7

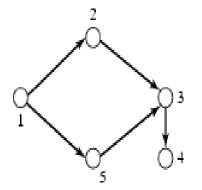
SELECT source FROM FallsInto CONNECT BY PRIOR source = dest WHERE Level <= 2 START WITH dest =1





## **Concept of Transitive Closure**

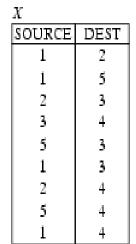
- •.Consider a graph G = (V, E)
- Let G\* = Transitive closure of G
- Then  $T = graph (V^*, E^*)$ , where
  - V\* = V
  - (A, B) in E\* if and only if there is a path from A to B in G.



R				
SOURCE	DEST			
1	2			
1	5			
2	3			
3	4			
5	3			

(a) Graph G

(b) Relation form



•Example:

- G has 5 nodes and 5 edges
- G\* has 5 nodes and 9 edges
- Note edge (1,4) in G\* for
  path (1, 2, 3, 4) in G.

(c) Transitive closure (G) = Graph  $G^*$ 

(d) Transitive closure in relational form 18



## 6.3.3 SQL3 Recursion

## •Computing table X from table R (figure in previous slide) WITH RECURSIVE X(source,dest)

AS (SELECT source,dest FROM R)

UNION

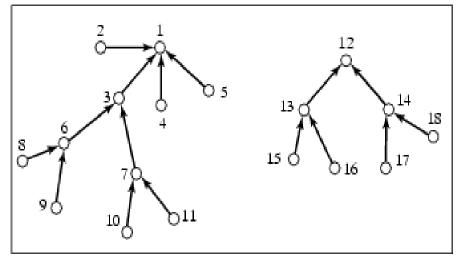
(SELECT R.source,X.dest FROM R,X WHERE R.dest=X.source);

- Meaning
  - Initialize X by copying directed edges in relation R
  - Infer new edge(a,c) if edges (a,b) and (b,c) are in X
  - Declarative query does not specify algorithm needed to implement it
- The graph can contain cycles!



With Recursive C(source, dest) as (Select source, dest From FallsInto Union Select FallsInto.source, C.dest From FallsInto, C Where FallsInto.dest = C.source) Select \* From C Where dest = 1;

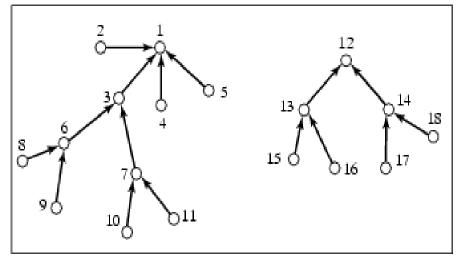
•Visit 2, 3, 4, 5, 6, 7, 8, 9, 10, 11





With Recursive C(source, dest) as (Select source, dest From FallsInto Union Select FallsInto.source, C.dest From FallsInto, C Where FallsInto.dest = C.source) Select \* From C Where source = 3;

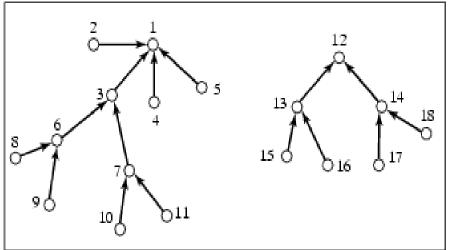
•Visit 6, 3, 1





With Recursive C(source, dest, depth) as (Select source, dest, 1 From FallsInto Union Select FallsInto.source, C.dest, C.depth+1 From FallsInto, C Where FallsInto.dest = C.source) Select \* From C Where dest = 1 And depth <= 2;

•Visit 2, 3, 4, 5, 6, 7





With Recursive C(source, dest, path) as

(Select source, dest, ARRAY[FallsInto.dest] From FallsInto Union All

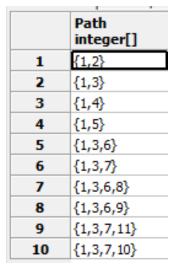
Select FallsInto.source, C.dest, path || FallsInto.dest

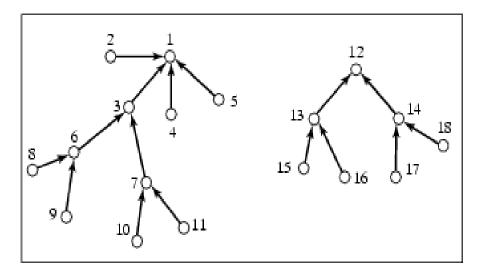
From FallsInto, C

**Where** FallsInto.dest = C.source)

Select path || source as "Path" From C

Where dest = 1;







## **Cycles**

With Recursive C(source, dest) as (Select source, dest From g Union Select g.source, C.dest From g, C Where g.dest = C.source) Select \* From C Where source = 1;

•Visit 2, 3, 4, 1

# $1 \xrightarrow{2} 3$

#### But be careful with Union All